

Peter Smith, *Introduction to Formal Logic* (CUP, 2nd edition)

## Exercises 35: Q-valuations

Thanks to “spamegg” (<https://github.com/spamegg1/>) for draft solutions!

(a) *Take a QL language, with a couple of proper names, a couple of unary predicates, and a binary predicate, and consider the following q-valuation:*

The domain is {Romeo, Juliet, Benedick, Beatrice}  
m: Romeo  
n: Juliet  
F: {Romeo, Benedick}  
G: {Juliet, Beatrice}  
L: {⟨Romeo, Juliet⟩, ⟨Juliet, Romeo⟩, ⟨Benedick, Beatrice⟩,  
⟨Beatrice, Benedick⟩, ⟨Benedick, Benedick⟩}.

*Then what are the truth values of the following wffs?*

(1)  $\exists xLmx$

True, intuitively: Romeo loves someone.

True, formally. According to our rule for evaluating existential wffs,  $\exists xLmx$  is true on the given valuation  $q$  just if there is an expanded valuation  $q_a$ , assigning some object in the domain as reference to the dummy name  $a$ , which makes  $Lma$  true. Which there is – take the expansion which assigns Juliet as reference to  $a$ .

(2)  $\forall xLxm$

False, intuitively: not everyone loves Romeo.

False, formally. According to our rule for evaluating universal wffs,  $\forall xLxm$  is false on the given valuation  $q$  just if there is an expanded valuation  $q_a$ , assigning some object in the domain as reference to the dummy name  $a$ , which makes  $Lam$  false. Which there is – take the expansion which assigns Benedick as reference to  $a$ .

(3)  $(\exists xLmx \rightarrow Lmn)$

True. A material conditional with a true antecedent (as we’ve shown), and a true consequent (obviously, since the pair ⟨Romeo, Juliet⟩ is in the extension of  $L$ ).

(4)  $\forall x(Fx \rightarrow \neg Gx)$

True intuitively. Putting it sloppily, the  $F$ s are the men in the domain, the  $G$ s are the women, and none of the men are women.

True, formally. According to our rule for evaluating universal wffs,  $\forall x(Fx \rightarrow \neg Gx)$  is true on the given valuation  $q$  just if *every* expanded valuation  $q_a$  makes  $(Fa \rightarrow \neg Ga)$ . Which it does. When  $a$  is assigned Romeo or Benedick, the antecedent and consequent of the conditional are both true; when  $a$  is assigned Juliet or Beatrice, the antecedent and consequent of the conditional are both false.

(5)  $\forall x(Gx \rightarrow (Lxm \vee \neg Lmx))$

True intuitively. Any woman in the domain is either loved by Romeo or doesn’t love him.

True, formally.  $\forall x(Gx \rightarrow (Lxm \vee \neg Lmx))$  is true on the given valuation  $q$  just if *every* expanded valuation  $q_a$  makes  $(Ga \rightarrow (Lam \vee \neg Lma))$ . Which it does. When  $a$  is assigned Romeo or Benedick, the antecedent of this conditional is false, so the whole conditional is true; when  $a$  is assigned Juliet,  $Ga$  and  $Lam$  are true, making the whole conditional true; when  $a$  is assigned Beatrice,  $Ga$  and  $\neg Lma$  are true, making the whole conditional true again.

(6)  $\forall x(Gx \rightarrow \exists yLxy)$

True intuitively: every woman in the domain loves someone.

True, formally. Because  $\forall x(Gx \rightarrow \exists yLxy)$  is true on the given valuation  $q$  just if *every* expanded valuation  $q_a$  makes (\*)  $(Ga \rightarrow \exists yLay)$  true. Which it does:

- i. When  $a$  is assigned Romeo or Benedick, the antecedent of this conditional is false, so (\*) is true.
- ii. When  $a$  is assigned Juliet, the antecedent of (\*) is true, and the consequent is true if  $\exists yLay$  is true on  $q_a$ , i.e. if there is some expanded valuation  $q_{ab}$  which makes  $Lab$  true – which there is, because we can assign Romeo to  $b$ .
- iii. Similarly when  $a$  is assigned Beatrice, the antecedent of (\*) is true, and the consequent is true if  $\exists yLay$  is true on  $q_a$ , i.e. if there is some expanded valuation  $q_{ab}$  which makes  $Lab$  true – which there is, because we can this time assign Benedick to  $b$ .

(7)  $\exists x(Fx \wedge \forall y(Gy \rightarrow Lxy))$

False intuitively: it's not the case that, in this domain, some man loves every woman.

Formally,  $\exists x(Fx \wedge \forall y(Gy \rightarrow Lxy))$  is true on  $q$  just if *some* expanded valuation  $q_a$  makes (\*)  $(Fa \wedge \forall y(Gy \rightarrow Laxy))$  true.

- i. But obviously an expanded valuation that assigns Juliet or Beatrice to  $a$  and hence makes  $Fa$  false can't make (\*) true.
- ii. What about the valuation which assigns Romeo to  $a$ ? Then (\*) is true just if  $\forall y(Gy \rightarrow Lay)$  is true on  $q_a$ , i.e. just if (\*\*)  $(Gb \rightarrow Lab)$  is true on *every* further expanded valuation  $q_{ab}$ . But assign Beatrice of  $b$ , and (\*\*) is false. So (\*) isn't true on  $q_a$  when  $a$  is assigned Romeo.
- iii. Similarly when  $a$  is assigned Benedick.

So (\*) comes on true on no expanded valuation  $q_a$ , so the wff (7) is false on  $q$ .

(b) *Now take the following  $q$ -valuation:*

The domain is  $\{4, 7, 8, 11, 12\}$

$m$ : 7

$n$ : 12

$F$ : the set of even numbers in the domain

$G$ : the set of odd numbers in the domain

$L$ : the set of pairs  $\langle m, n \rangle$  where  $m$  and  $n$  are in the domain and  $m < n$ .

*What are the truth values of the wffs (1) to (7) now?*

(1) True

(2) False

(3) True

The antecedent is true by (1), the consequent  $Lmn$  is true because  $7 < 12$ .

(4) True

(5) False.

$\forall x(Gx \rightarrow (Lxm \vee \neg Lmx))$  is true on the given valuation  $q$  just if *every* expanded valuation  $q_a$  makes (\*)  $(Ga \rightarrow (Lam \vee \neg Lma))$  true.

But take the valuation which assigns 11 to  $a$ . Then  $Ga$  is true,  $Lam$  is false and  $\neg Lma$  is false. Which makes (\*) false.

(6) True.

In our domain, every odd number is less than some number.

(7) True.

In our domain, there is an even number (4, in fact) which is less than every odd number.

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(c) Take the language  $QL_3$  of Exercises 30(b) whose non-logical vocabulary comprises the name  $n$ , the one-place predicates  $F, G, H$ , the two-place  $L$ , and three-place predicate  $R$ . Consider the following  $q$ -valuation (the natural one suggested by the given glossary for the language):

The domain is the set of natural numbers (the integers from zero up)

$n$ : one

$F$ : the set of odd numbers

$G$ : the set of even numbers

$H$ : the set of prime numbers

$L$ : the set of pairs  $\langle m, n \rangle$  such that  $m < n$

$R$ : the set of triples  $\langle l, m, n \rangle$  such that  $l = m + n$ .

Carefully work out the values of the wffs (1), (2), (4) and (5) from Exercises 30(b), i.e. the wffs

(1)  $\forall x \forall y \exists z Rzxy$

True.

Informally this says, given any two numbers, they have a sum. (“The natural numbers are closed under addition.”)

Formally, (1) is true on the given valuation  $q$  if for every expansion  $q_{ab}$  – i.e. whatever values we assign  $a$  and  $b$  – (\*)  $\exists z Rzab$  is true. Which it is!

(2)  $\exists y \forall x Lxy$

False.

Informally it translates to: “there is a number bigger than all others.”

Formally, we could argue by contradiction. So assume (2) evaluates as true on the given  $q$ -valuation  $q$ . Then there exists an expanded  $q$ -valuation  $q_a$ , where the dummy name  $a$  is assigned the natural number  $a$ , such that  $\forall x Lxa$  is true.

That means that on every further expanded  $q$ -valuation  $q_{ab}$ , whatever  $b$  is assigned  $Lba$  is true on  $q_{ab}$ . But that’s absurd. Suppose  $b$  is assigned  $a + 1$ !

(4)  $\forall x (Hx \rightarrow \exists y (Lxy \wedge Hy))$

This is true on the given valuation  $q$  is true iff for any extended valuation  $q_a$  (where  $a$  now denotes the number  $a$ ) the wff (\*)  $(Ha \rightarrow \exists y (Lay \wedge Hy))$  is true.

But the conditional (\*) is true on any  $q_a$  because its consequent is true on any  $q_a$  – there is always an extension  $q_{ab}$  (where  $b$  now denotes the number  $b$ ) which makes  $Lab \wedge Hb$  true. Because for any  $a$  there is a larger prime  $b$ .

(There a famous and elegant ancient Greek proof from Euclid’s *Elements* for any number there is a larger prime. You can look it up on the web if you don’t know it!)

(5)  $\forall x \forall y ((Fx \wedge Ryxn) \rightarrow \neg Fy)$

True if and only if on every extended valuation  $q_{ab}$  (whatever those dummy names  $a, b$  pick out) the wff  $((Fa \wedge Rban) \rightarrow \neg Fb)$  is true. Which it is, because if  $a$  is odd, and  $b = a + 1$ , then  $b$  is not odd.

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In the book there is a final exercise (d\*): *Show that if the wff  $\alpha$  doesn't contain the dummy name  $\delta$ , then  $\alpha$  is true on the valuation  $q$  if and only if it is also true on any expansion  $q_\delta$ .*

Obviously, the idea is that if  $\alpha$  *doesn't* contain the dummy name  $\delta$  then what value we give  $\delta$  on a valuation can't affect the value of  $\alpha$ .

And that's basically right, but the devil is the details. I have to confess a slight mess-up here. For example, how would the story go for the value of  $\forall xFx$  on the extended valuation  $q_a$ ?

Remember we say that  $\forall xFx$  is true on a valuation if  $Fa$  (plugging in the first available dummy name) is true on every expansion of that valuation assigning a value to  $a$ , whatever value we choose. So, applied to the present case,  $\forall xFx$  is true on the extended valuation  $q_a$  if it is true on every expansion of *that* valuation which assigns a value to  $a$ . But hold on! By definition  $q_a$  already assigns an object to  $a$ . 'Expansions' here is the wrong word – to get what we want, we must here take ourselves to be considered variations of the valuation  $q_a$  which spin the interpretation of  $a$ . But this doesn't really tally with the strict letter of what is in the book! Oooops.

This glitch hopefully doesn't affect what happens elsewhere in the book. We could fiddle with the definition of a  $\delta$ -expansion at the top of p. 35. But in fact, looking at this chapter, I think some rewriting is really needed both to tidy up the glitch but also to make things clearer overall.