

Exercises 29: Simple QL translations

In the language QL_1 , the *proper names* are just

m: Socrates,
n: Plato,
o: Aristotle;

and the *predicates* are just

F: ① is a philosopher,
G: ① is a logician,
H: ① is wise,
L: ① loves ②,
M: ① is a pupil of ②,
R: ① prefers ② to ③.

The *domain of quantification* for QL_1 : people, past and present.

(a) Translate the following into QL_1 as best you can:

- (1) Aristotle is a pupil of Socrates and Plato.
- (2) Plato taught Aristotle only if Socrates taught Plato.
- (3) If Plato is a pupil of someone, he is a pupil of Socrates.
- (4) Some philosophers are not wise.
- (5) Not all philosophers are wise.
- (6) Any logician loves Aristotle.
- (7) No one who is a wise philosopher prefers Plato to Aristotle.
- (8) Whoever is a pupil of Plato is wise.
- (9) Not every wise logician is a pupil of Aristotle.
- (10) Any logician is a wise philosopher.
- (11) Aristotle prefers no philosopher to Plato.
- (12) Some wise people aren't philosophers, and some aren't logicians.
- (13) Only philosophers love Aristotle.
- (14) Not only philosophers love Socrates.
- (15) Socrates is a philosopher whom everyone wise loves.
- (16) Only some logicians love Plato.
- (17) No philosophers or logicians are wise.
- (18) All philosophers and logicians are wise.

(b) Which of these pairs of wffs are equivalent, which not, and why?

- (1) $(\exists xFx \vee \exists xGx), \exists x(Fx \vee Gx)$
- (2) $(\exists xFx \wedge \exists xGx), \exists x(Fx \wedge Gx)$
- (3) $(\exists xFx \wedge \exists xGx), \exists x(Fx \wedge \exists yGy)$
- (4) $(\exists xFx \wedge \exists xGx), \exists x\exists y(Fx \wedge Gy)$
- (5) $\exists x\forall y(Fx \wedge Gy), \forall y\exists x(Fx \wedge Gy)$
- (6) $\exists x\forall y(Fx \vee Gy), \forall x\exists y(Fy \vee Gx)$
- (7) $(\forall xFx \rightarrow \forall xGx), \exists x(Fx \rightarrow \forall yGy)$
- (8) $(\forall xFx \rightarrow \forall xGx), \exists x\forall y(Fx \rightarrow Gy)$
- (9) $(\forall xFx \rightarrow \forall xGx), \forall y\exists x(Fx \rightarrow Gy)$

(c*) Use equivalences you now know about to outline a proof that every wff is equivalent to one in *prenex* form, where all quantifiers are at the beginning of the wff.

(d*) We can render 'Plato and Aristotle are philosophers' by e.g. $(Fn \wedge Fo)$. Why can't we render

‘Plato and Aristotle are classmates’ by e.g. $(C_n \wedge C_o)$? What does this sort of case tell us about some expressive limitations of QL languages?