

Exercises 18: The truth-functional conditional

(a) Suppose we are working in a PL language where ‘P’ means *Putnam is a philosopher*, ‘Q’ means *Quine is a philosopher*, etc. Translate the following as best you can:

- (1) If either Quine or Putnam is a philosopher, so is Russell.
- (2) Only if Putnam is a philosopher is Russell one too.
- (3) Quine and Russell are both philosophers only if Sellars is.
- (4) Russell’s being a philosopher is a necessary condition for Quine’s being one.
- (5) Russell’s being a philosopher is a sufficient condition for Quine’s being one.
- (6) Putnam is a philosopher if and only if Quine isn’t.
- (7) Provided that Quine is a philosopher, Russell is one too.
- (8) Quine is not a philosopher unless Russell is one.
- (9) Only if either Putnam or Russell is a philosopher are both Quine and Sellars philosophers.

(b) Assuming that we are dealing with a suitable PL language. Which of the following arguments *ought* to come out valid, assuming that ‘ \rightarrow ’ is a reasonably good surrogate for ‘if ..., then ...’? Which is tautologically valid?

- (1) $P, (P \rightarrow Q), (Q \rightarrow R) \therefore R$
- (2) $\neg R, (P \rightarrow R), (Q \rightarrow P) \therefore \neg Q$
- (3) $(P \rightarrow \neg(Q \vee R)), (Q \rightarrow R), (\neg R \rightarrow P) \therefore (P \rightarrow R)$
- (4) $(P \vee Q), (P \rightarrow R), \neg(Q \wedge \neg R) \therefore R$
- (5) $(R \rightarrow (\neg P \vee Q)), (P \wedge \neg R) \therefore \neg(\neg R \vee Q)$
- (6) $(\neg P \vee Q), \neg(Q \wedge \neg R) \therefore (P \rightarrow R)$
- (7) $(P \wedge \neg R), (Q \rightarrow R) \therefore \neg(P \rightarrow Q)$
- (8) $\neg(\neg S \rightarrow (\neg Q \wedge R)), (P \vee \neg\neg Q), (R \vee (S \rightarrow P)) \therefore (P \rightarrow S)$

(c) Which of the following are true for all α, β, γ in a PL language and why? Which of the true claims correspond to true claims about the vernacular (bi)conditional?

- (1) If $\alpha, \beta \models \gamma$ then $\alpha \models (\beta \rightarrow \gamma)$.
- (2) $((\alpha \wedge \beta) \rightarrow \gamma) \models (\alpha \rightarrow (\beta \rightarrow \gamma))$.
- (3) $((\alpha \vee \beta) \rightarrow \gamma) \models ((\alpha \rightarrow \gamma) \vee (\beta \rightarrow \gamma))$.
- (4) If $\models (\alpha \rightarrow \beta)$ and $\models (\beta \rightarrow \gamma)$, then $\models (\alpha \rightarrow \gamma)$.
- (5) If $\models (\alpha \rightarrow \beta)$ and $\models (\alpha \rightarrow \neg\beta)$, then $\models \neg\alpha$.
- (6) $\models (\alpha \leftrightarrow \alpha)$
- (7) $(\alpha \leftrightarrow \beta) \models (\beta \leftrightarrow \alpha)$.
- (8) $(\alpha \leftrightarrow \beta), (\beta \leftrightarrow \gamma) \models (\alpha \leftrightarrow \gamma)$.
- (9) If $\models \alpha \leftrightarrow \beta$ then α and β are tautologically consistent.
- (10) If $\models \alpha \leftrightarrow \neg\beta$ then α and β are tautologically inconsistent.

(d*) On alternative languages for propositional logic:

- (1) Suppose the language PL_1 has just the connectives \rightarrow and \neg (with the same interpretation as before). Show that disjunction and conjunction can be expressed in PL_1 . Conclude that PL_1 has an expressively adequate set of built-in connectives.
- (2) Consider too the variant language PL_2 whose only logical constants are \rightarrow and the absurdity constant \perp . Show that in PL_2 we can introduce a negation connective so that $\neg\alpha$ is shorthand for $(\alpha \rightarrow \perp)$. Conclude that PL_2 is also expressively adequate.